## FINAL CLASS!

My next job?



## Announcements

- Review Session Helpful? Solutions out soon.
- Request Topics review.
- A single handwritten sheet of notes allowed for finals.
(Confirmation on this soon.)
- No electronic devices.


## Today?

## FINSTEDOUTHENMEP

## CS 354 ~ What to take away?

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## General

Everything is represented as a sequence of bits (0s and 1s)!

- Your executable
- Your images
- Your browser application
- Your pdfs
- Everything!

What they actually represent...

- Depends on context
- How we choose to interpret them.


## General

More specifically,

- I have a piece of data that is $\mathbf{0 1 0 0} \mathbf{0 0 0 1}$. What is it?
- What is 010010101011011010111010101010101010100 ?

Are we going to split it into 8 bit groups? Or split it into 32 bit groups?

Special Mention: Quantum Computing

## Compilation System

|  |  |  |  |  |  | printf.o |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| hello.c | Preprocessor (cpp) | hello.i | Compiler (cc1) | hello.s | $\begin{gathered} \text { Assembler } \\ \text { (as) } \end{gathered}$ | hello.o | Linker <br> (ld) | hello |
| Source |  | Modified source |  | Assembly |  | Relocatable |  | Executable object |
| program (text) |  |  |  | program <br> (text) |  | object programs |  |  |
|  |  | (text) |  |  |  | (binary) |  |  |

## C - Pointers

## Operators

```
&varX - ADDRESS OF variable varX.
*varY - VALUE AT address varY. (Indirection operator)
```

How do you define a char pointer?
char ch = 'y';
char * ptr = \&ch;

This right here is not an indirection operator. It is just pointer declaration syntax;

## Data Representation

## Endianness

Integer - 0x12 345678 - 4 bytes.

Little Endian

| Addr | Data |
| :---: | :---: |
| $0 \times 100$ | 78 |
| $0 \times 101$ | 56 |
| $0 \times 102$ | 34 |
| $0 \times 103$ | 12 |

Big Endian

| Addr | Data |
| :---: | :---: |
| $0 \times 100$ | 12 |
| $0 \times 101$ | 34 |
| $0 \times 102$ | 56 |
| $0 \times 103$ | 78 |

## Data Representation

4 bit datatype - 0000 to 1111

## Unsigned Representation

No specific bits to denote sign. So value goes from DEC 0 to 15.

## Signed Representation

MSB allocated for sign.
So value ranges from

| BIN | 0111 | 1000 |
| :---: | :---: | :---: |
| DEC | 7 | -8 |

## Assembly

- How to read the x86 Instruction Sheet.
- Registers
- Control Flags and Conditional Jumps.
- cmp and test instructions followed by jumps.
- Function Stack Frames!
- All the space associated with a function goes away after it returns i.
e. Popping it off the stack.


## Memory

- Not all storage technology are created equal.
- Some are fast and expensive and some are relatively slow and inexpensive.
- Use a fast memory to serve as a staging area for data from a slower one - CACHING!


## Memory Hierarchy



## Cache



Design a cache for me!
(S, E, B, m)

## Locality

## Temporal Locality

Accessing the same memory location over and over againGood Thing!

## Spatial Locality

Accessing the memory in sequence (stride-1 reference pattern) is a good thing as well.

Tip: Algorithm design does not consider the physical limitations of memory.

## Virtual Memory

## Why Virtual Memory?

- Memory Protection
- Easy Memory Management
- Use more space that what is physically available in the physical DRAM memory.


## Virtual Memory

## Important Task

$$
\text { Virtual Address } \longrightarrow \text { Physical Address }
$$

How?

- Memory Management Unit (In-charge of doing this!)
- Page Tables (Software entity)
- Translation Lookaside Buffer (Separate cache)


## Dynamic Memory Allocation

Huge space of unallocated memory - The Heap!
Why an allocator?

- Like any shared resource, access to this resource has to be controlled.
- See Tragedy of the Commons. (Shared resource - Road.

Controller - Cops and Traffic Rules).

- So, we need a Dynamic Memory Allocator.
- ... with a lot of rules and control mechanisms - Headers, Alignment Restrictions, Maintenance.


## Exceptions

Abrupt change in a processor's control flow.

## Different Kinds

- Interrupts (not under our control ~ Asynchronous)
- Press Ctrl + C!
- System Calls (we cause the abrupt change)
- Write to display..
- Faults (we cause this)
- Accessing a page that is not in the main memory.
- Abort (we cause this)


## Context Switching

Switch from executing one process to another... while storing the status (or context) of the switched out process.

## Why?

CPU should not wait for a slow task that a process needs to be performed. (Relative times - 1 sec for a CPU is 10 months for a Hard Disk)

Process has two modes - Kernel and User
Needs to be in kernel mode before performing a context switch. Why?

## Linking

Hope this is still fresh in memory!

## Goodbye

\&
Good Luck!

